

Forking and Dividing  
at a  
Generic Scale

An Introduction to Kim-Independence  
and  $NSOP_1$

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# Lecture 2

Last time

## Definition

$\Psi(x; y)$  has SOP, if there is

$(a_\eta)_{\eta \in 2^{<\omega}}$  such that

- (paths are consistent)  
For all  $\eta \in 2^{<\omega}$

$$\{ \Psi(x; a_{\eta \upharpoonright i}) : i < \omega \}$$

is consistent.

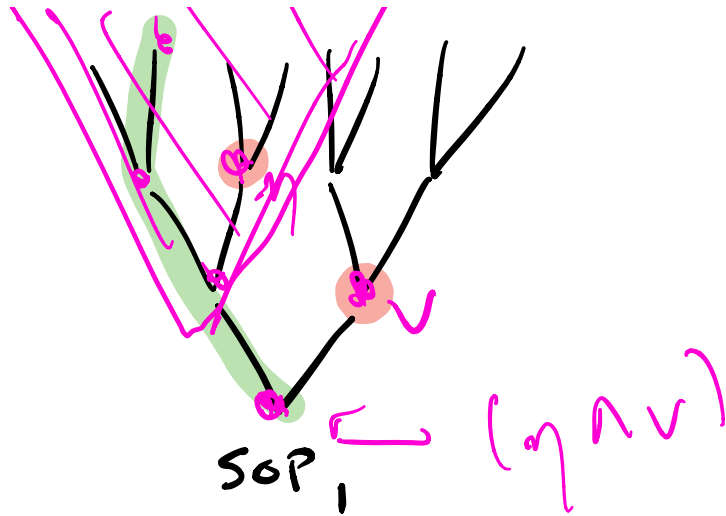
- (Inconsistency)

For all  $\eta, \nu \in 2^{<\omega}$ , if

$$\eta \sqsupseteq (\eta \cap \nu) \hat{\ } \langle 0 \rangle, \nu = (\eta \cap \nu) \hat{\ } \langle 1 \rangle$$

$$\{ \Psi(x; a_\eta), \Psi(x; a_\nu) \}$$

is inconsistent.



We say  $T$  has  $SOP_1$  if some formula has  $SOP_1$  modulo  $T$ ; we say  $T$  is  $NSOP_1$  if it's not  $SOP_1$ .

# SOP, Arrays

## Proposition

The following are equivalent:

(1)  $T$  has SOP,

(2) There is  $\varphi(x; y)$ ,  $k \in \omega$ , and

$(c_{i,0}, c_{i,1})_{i \in \omega}$  such that

(a)  $\{\varphi(x; c_{i,0}) : i \in \omega\}$  is consistent.

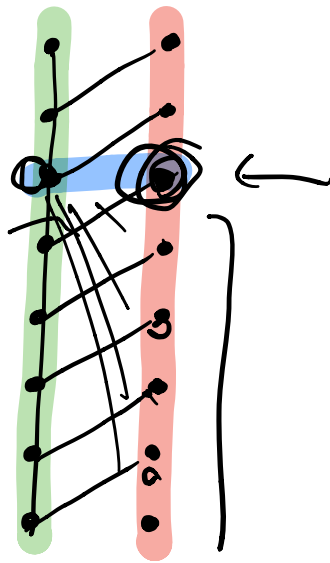
(b)  $\{\varphi(x; c_{i,1}) : i \in \omega\}$  is  $k$ -inconsistent.

(c) For all  $i \in \omega$ ,  $c_{i,0} \equiv_{c_{i,0} c_{i,1}}$   $c_{i,1}$

(3) Same as (2), but with

$k=2$ .





SOP<sub>i</sub>

Note The equivalence of (1) and (3) also holds for formulas. Also, since any  $\mathcal{U}$  will witness SOP<sub>i</sub> also in any expansion of our ambient monster  $\mathcal{M}$  (in the language  $L$ , say), we may find such an array with

$$c_{i,0} \equiv_{c_{i,0}, c_{i,1}}^{L'} c_{i,1} \text{ where } L' \text{ is a larger language,}$$

once  $M$  has been expanded to an  $L'$ -structure  $M'$ .

## Invariant Types

Recall  $p(x) \in S_x(M)$  is called

A-invariant if  $\sigma p = p$  for all  $\sigma \in \text{Aut}(M/A)$

or, equivalently, for all  $\varphi(x; y)$ , if

$c \equiv_A c'$ , then

$$\varphi(x; c) \in p \iff \varphi(x; c') \in p$$

Given A-invariant  $p \in S_x(M)$ ,  $q \in S_y(M)$ ,

we obtain an A-invariant type

$(p \otimes q)(x, y) \in S_{x, y}(M)$  defined by

$$(p \otimes q) = \text{tp}(a, b / M)$$

for some  $b \models q|_M$  and  $a \models p|_{M_b}$ .

Given any linear order  $I$  and

$A$ -invariant  $p \in S_x(M)$ , we have

another  $A$ -invariant global type

$p^{\otimes I}(x_\alpha : \alpha \in I)$  such that

$$(a_\alpha)_{\alpha \in I} \models p^{\otimes I} \Leftrightarrow a_\alpha \models p|_{M a_\alpha} \text{ for all } \alpha \in I.$$

If  $B \geq A$  and  $(a_\alpha)_{\alpha \in I} \models p^{\otimes I}|_B$ ,

then  $(a_\alpha)_{\alpha \in I}$  is  $B$ -indiscernible.

When  $I$  is an infinite linear order,

we refer to  $(a_\alpha)_{\alpha \in I}$  as a

Morley sequence in  $p$  (over  $B$ ).

Note: Being a Morley sequence over  $A$  in a global  $A$ -invariant type  $p$  implies being a Morley sequence over  $A$  in the sense of the previous lecture, i.e.

$$(a_\alpha)_{\alpha \in I} \text{ f.p. } \mathbb{F}^{\text{op}} \big|_A \Rightarrow a_\alpha \downarrow_A^f a_{\alpha'} \text{ for all } \alpha \neq \alpha' \in I.$$

But the converse is usually false.

## Average Types

If  $A$  is a set of parameters and  $\mathcal{D}$  is an ultrafilter on  $A^n$ , then for any  $B$ ,

$$A_v(D, \mathcal{B}) = \left\{ \varphi(x) \in L(\mathcal{B}) \mid \{a \in A \mid \models \varphi(a)\} \in D \right\}$$

is called the average type of  $D$  over

$\mathcal{B}$ .

(1)  $A_v(D, \mathcal{B})$  is a complete type over  $\mathcal{B}$ .

(2)  $A_v(D, M)$  is a global  $A$ -invariant type (in fact,  $A$ -finitely satisfiable).

(3) If  $M \models T$  and  $p \in S_x(M)$ , there is some  $D$  on  $M^{|x|}$  such that  $A_v(D, M) = p$ .

Together, these imply the following

## Fact

"Types over models have global invariant extensions"

If  $M \models T$  and  $p \in S_x(M)$ , then there is a global  $M$ -invariant (even  $M$ -finitely satisfiable)  $q \in S_x(M)$  with  $q|_M = p$ .

==

Proof Pick an ultrafilter  $\mathcal{D}$  on  $M^{|x|}$  such that  $A_v(\mathcal{D}, M) = p$ .

Then  $A_v(\mathcal{D}, M)$  is the desired global extension.

## Note

This fact implies that, over models, in every type, we have some notion of a 'generic sequence'—namely, a 'generic sequence in  $p'$ ', for some  $p \in S_x(M)$ , should be a Morley sequence over  $M$  in some global  $M$ -invariant type extending  $p$ .

# Kim-Independence

## Definition

Fix  $M \models T$ .

(1) We say  $\varphi(x/a)$  Kim-divides over  $M$  if there is a global  $M$ -invariant  $p \equiv \text{tp}(a/M)$  such that  $\varphi$  divides with respect to a Morley sequence over  $M$  in  $p$  (or, equivalently, in every Morley sequence over  $M$  in  $p$ ).



(2) We say  $\psi(x; a)$  Kim-forks over  $M$

if  $\psi(x; a) \vdash \bigvee_{i \leq k} \psi_i(x; c_i)$ , where

$\psi_i(x; c_i)$  Kim-divides over  $M$ .

(3) We say a type  $p(x)$  Kim-divides

or Kim-forks over  $M$  if  $p$  implies

some formula that Kim-divides

or Kim-forks over  $M$ , respectively.

(4) We write  $a \underset{M}{\downarrow}^k b$  to

indicate  $\text{tp}(a/Mb)$  does not

Kim-fork over  $M$ .

## Kim's Lemma

### Theorem (Kaplan-R.)

The following are equivalent:

(1)  $T$  is NSOP<sub>1</sub>

(2) Kim's lemma for Kim-dividing

If  $M \models T$  and  $\mathcal{U}(x; a)$

Kim-divides over  $M$ , then

$\mathcal{U}(x; a)$  divides with respect to

Morley sequences over  $M$  in

every global  $M$ -invariant type

$p \geq \text{tp}(a/M)$ .

(3) If  $M \neq T$ , then if  $\varphi(x; a)$  divides with respect to a Morley sequence over  $M$  in some global  $M$ -finitely satisfiable type  $p \geq tp(a/M)$ , then  $\varphi(x; a)$  divides with respect to Morley sequences over  $M$  in every global  $M$ -finitely satisfiable  $q \geq tp(a/M)$ .

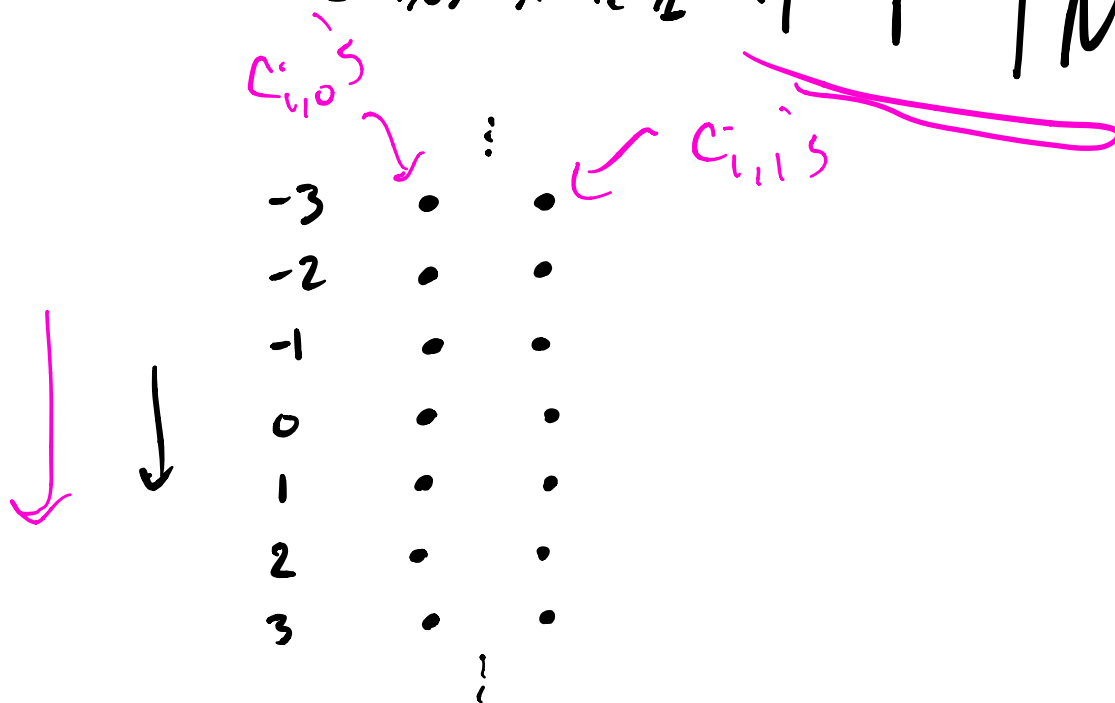
Proof  
(2)  $\Rightarrow$  (3) follows from the fact that  $M$ -finitely satisfiable types are  $M$ -invariant.

(1)  $\Rightarrow$  (2) Assume (2) fails and we will show  $T$  has SOP<sub>1</sub>.

So fix  $M \models T$ ,  $\varphi(x; a)$ , and global  $M$ -invariant  $p, q \geq \text{tp}(a/M)$  such that  $\varphi(x; a)$  divides with respect to Morley sequences over  $M$  in  $q$ , but  $\varphi(x; a)$  does not divide with respect to Morley sequences over

$M$  in  $\rho$ .

Now let  $(c_{i,0}, c_{i,1})_{i \in \mathbb{Z}} \models (p \otimes q)^{\otimes \mathbb{Z}} \mid M$ .



The rest of the proof will be a reflection on this picture.

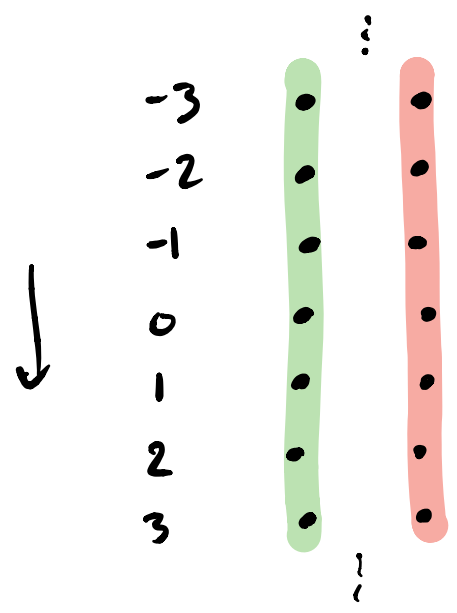
Observation 1  $(c_{i,0})_{i \in \mathbb{Z}} \models p^{\otimes \mathbb{Z}} \mid M$

so  $\{\varphi(x; c_{i,0}) : i \in \mathbb{Z}\}$  is consistent.



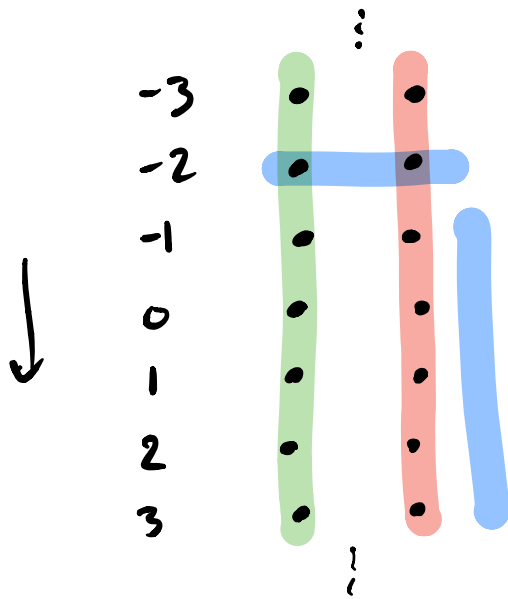
Observation 2 :  $(c_{1,i})_{i \in \mathbb{Z}} \stackrel{k}{\neq} q^{\otimes \mathbb{Z}} / M$

so  $\{\psi(x; c_{1,i}) : i \in \mathbb{Z}\}$  is  $k$ -inconsistent for some  $k$ .



Observation 3: For all  $j \in \mathbb{Z}$ ,

$$c_{j,0} \equiv M c_{j,1} c_{j,1} \leftarrow$$



Why? Notice

$$(c_{i,0}, c_{i,1})_{i \in \mathbb{Z}} \in (p \otimes q)^{\otimes \mathbb{Z}} / M$$

So

$$(c_{i,0}, c_{i,1})_{i \geq j} \in (p \otimes q)^{\otimes \omega} / \underbrace{M c_{j,0}, c_{j,1}}$$

Moreover,  $c_{j,0} \neq p|_M$ ,  $c_{j,1} \neq q|_M$  and

$$p|_M = q|_M = \text{tp}(a/M) \text{ so}$$

$$c_{j,0} \equiv_M c_{j,1}. \leftarrow$$

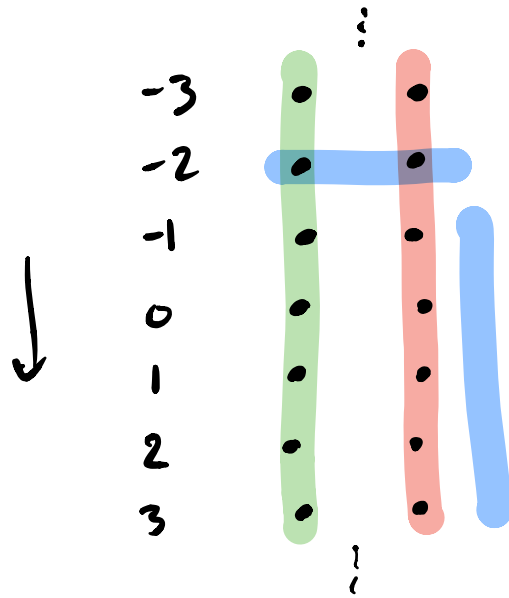
Because  $(p \otimes q)^{\text{ow}}$  is  $M$ -invariant,  
we know, for all  $\varphi(\bar{z}; y) \in L(M)$

$$\varphi(\bar{z}; c_{j,0}) \in (p \otimes q)^{\text{ow}} \Leftrightarrow \varphi(\bar{z}; c_{j,1}) \in (p \otimes q)^{\text{ow}}$$

so we have

$$c_{j,0} \equiv_{M_{c_{j,0} c_{j,1}}} c_{j,1} \leftarrow$$





So we've found an SOP, array,  
upside down. [To put it right side  
up, define  $(d_{i,0}, d_{i,1})_{i \in \omega}$  by

$$(d_{i,0}, d_{i,1}) = (c_{-i,0}, c_{-i,1})$$

for all  $i \in \omega$ .] This shows

$$\neg(2) \Rightarrow \neg(1) \quad \text{so} \quad (1) \Rightarrow (2).$$

(1)  $\Rightarrow$  (3) Assume  $T$  has  $SOP_2$ .

Let  $T^{Sk}$  be an expansion of  $T$  to a larger language that contains Skolem functions — i.e. such that for every  $\varphi(x; y)$  in  $L^{Sk}$  with  $|x|=1$  there is a function  $f_\varphi(y)$  such that

$$T^{Sk} \vdash (\forall y) \left[ (\exists x) \varphi(x; y) \rightarrow \varphi(f_\varphi(y); y) \right].$$

Then there is an  $L$ -formula  $\varphi(x; y)$  and a collection of types  $(c_{i,0}, c_{i,1})_{i \in \omega}$  such that

- $\{\varphi(x; c_{i,0}) : i < \omega\}$  is consistent.
- $\{\varphi(x; c_{i,1}) : i < \omega\}$  is 2-inconsistent.
- $c_{i,0} \equiv_{L^{sk}} c_{i,1}$  for all  $i < \omega$ .

By Ramsey and compactness, we may assume further that

$\langle (c_{i,0}, c_{i,1}) : i < \omega \rangle$  is an  $L^{sk}$ -indiscernible sequence.

Let  $M = \text{dcl}_{L^{sk}}(\langle c_{i,0}, c_{i,1} \rangle_{i < \omega})$ .

Then  $M \upharpoonright L \models T$ .

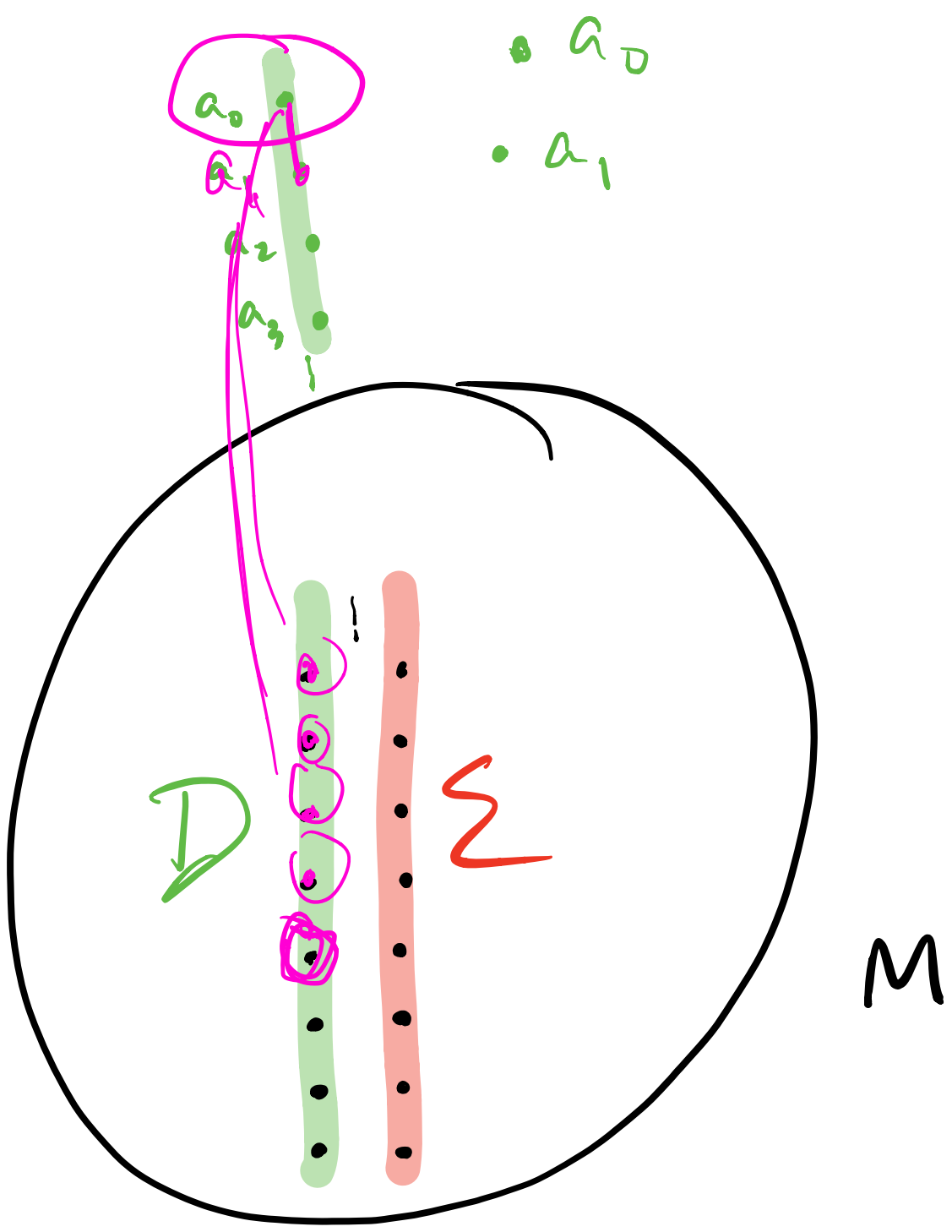
Let  $n = |c_{i,0}| = |c_{i,1}|$ , and let  $\mathcal{D}$  be a non-principal ultrafilter on  $M^n$ , concentrating on  $\{c_{i,0} : i < \omega\} \subset M^n$

and let  $\mathcal{E}$  be a non-principal  
ultrafilter concentrating on  
 $\{c_i : i < \omega\}$ .

Let  $p = A_v(D, M)$   
 $q = A_v(\mathcal{E}, M)$ .

Observation 1: If  $(a_i)_{i < \omega}$   
is a Morley sequence over  $M$  in  $p$   
then  $\{\varphi(x, a_i) : i < \omega\}$  is consistent.

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$\bullet a_0$   
 $\bullet a_1$

$a_0$

$a_1$

$a_2$

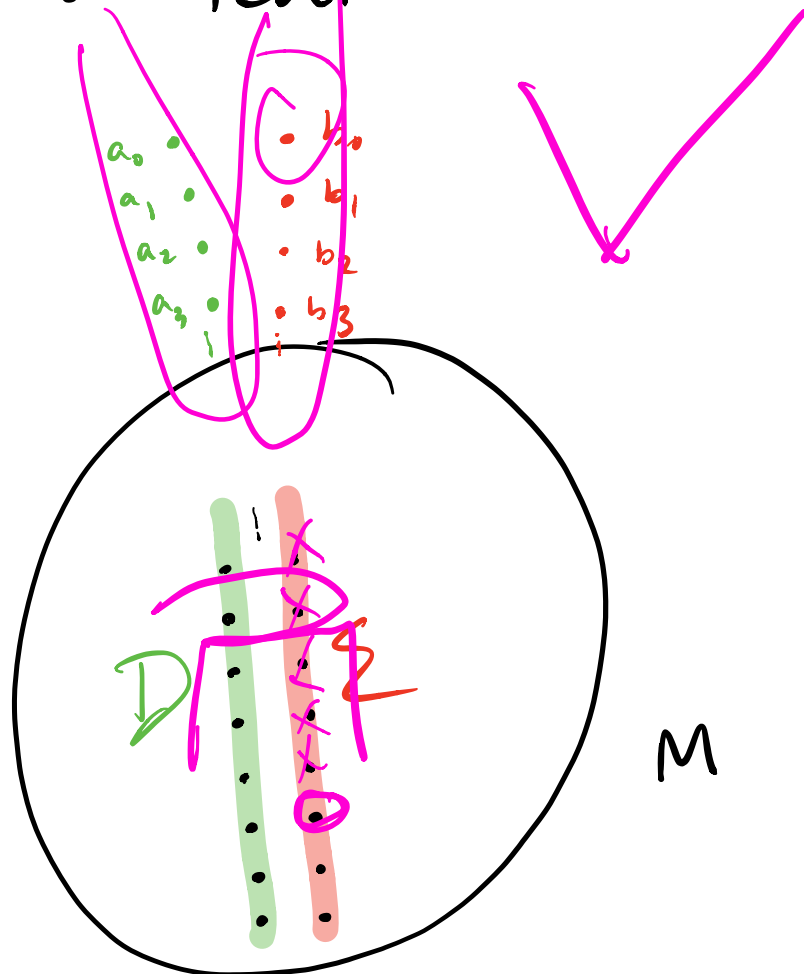
$a_3$

$D$

$\Sigma$

$M$

Observation 2: If  $(b_i)_{i \leq w}$  is a Morley sequence over  $M$  in  $\mathcal{L}$ , then  $\{\varphi(x; b_i) : i \leq w\}$  is 2-inconsistent.



Observation 3 :

$$P|M = Q|M.$$

Suppose  $\psi(x; m) \in L(M)$ . Then there is a tuple of  $L^{sk}$ -terms  $t$  and some  $k$  such that  $\psi(x; m)$  is equivalent modulo  $\top^{sk}$  to

$$\psi(x; t(c_{k,0}, c_{k,1})).$$

Because  $c_{i,0} \equiv_{L^{sk}}^{c_{i,0}, c_{i,1}} c_{i,1}$  for all  $i < \omega$ , we know, in particular

$$c_{j,0} \equiv_{L^{sk}}^{c_{k,0}, c_{k,1}} c_{j,1}$$

for all  $j > k$ . Hence

$$\psi(x_{i,m}) \in \mathcal{P} | \mathcal{M} \Leftrightarrow \{c_{j,0} \mid \nexists \psi(c_{j,0}, m)\} \in \mathcal{D}$$

$$\Leftrightarrow \{c_{j,0} \mid \nexists \psi(c_{j,0}; t(c_{k,0}, c_{k,1}))\} \in \mathcal{D}$$

$$\Leftrightarrow \{c_{j,0} \mid \nexists \psi(c_{j,0}; t(c_{k,0}, c_{k,1}))\} \supseteq \{c_{j,0} \mid j > k\}$$

$$\Leftrightarrow \{c_{j,1} \mid \nexists \psi(c_{j,1}; t(c_{k,0}, c_{k,1}))\} \supseteq \{c_{j,1} \mid j > k\}$$

$$\Leftrightarrow \{c_{j,1} \mid \nexists \psi(c_{j,1}, m)\} \in \mathcal{E}$$

$$\Leftrightarrow \psi(x_{i,m}) \in \mathcal{Q} | \mathcal{M}. \quad \#$$



Kim-forking = Kim-dividing

Proposition

Assume  $T$  is NSOP<sub>1</sub>. Then

if  $M \models T$  and  $\varphi(x; a)$  Kim-forks over  $M$ , then  $\varphi(x; a)$  Kim-divides over  $M$ .

Proof Fix  $M \models T$ .

Suppose  $\varphi(x; a) \vdash \bigvee_{i \leq k} \psi_i(x; c_i)$

where  $\psi_i(x; c_i)$  Kim-divides

over  $M$  for all  $i \leq k$ .

Let  $q \equiv \text{tp}(a, c_0, \dots, c_{k-1} / M)$

be a global  $M$ -invariant type,  
and let

$$\langle \underline{a_i}, \underline{c_{0,i}}, \underline{c_{1,i}}, \dots, \underline{c_{k-1,i}} \rangle \neq \varnothing^{\oplus \omega} / M.$$

Then  $(a_i)_{i \in \omega}$  and each  
 $(c_{j,i})_{i \in \omega}$  are also Morley  
sequences over  $M$  in global  
 $M$ -invariant types. By Kim's  
lemma for Kim-dividing, we  
know

$$\{\psi_j(x; c_{j,i}) : i < \omega\}$$

is inconsistent for all  $j < k$ .

We want to show

$$\{\varphi(x; a_i) : i \in \omega\}$$

is inconsistent.

Towards contradiction, suppose not  
and let  $b$  be a realization.

Then we know, for all  $i \in \omega$ ,

$$\varphi(x; a_i) \vdash \bigvee_{j < k} \varphi_j(x; c_{j,i})$$

So for all  $i \in \omega$ , there is some  
 $\bar{j}(i) < k$  such that

$$\models \varphi_{\bar{j}(i)}(b; c_{\bar{j}(i), i}).$$

By the Pigeonhole principle,  
there is some  $j_* < k$  such that  
the set  $X := \{i < \omega \mid j(i) = j_*\}$   
is infinite. Then we have  
both  $\{\mathcal{A}_{j_*}(x; c_{j_*, i}) : i \in X\}$   
is inconsistent (by indiscernibility)  
and that this set of formulas  
is realized by  $b$ , a  
contradiction.  $\square$

# The Kim-Pillay Criterion

## NSOP<sub>1</sub>

### Theorem (Chernikov-R.)

Suppose  $\perp^*$  is an  $\text{Aut}(M)$ -invariant ternary relation on small subsets of  $M$

satisfying, for all  $M \models T$ :

- (1) Existence:  $a \perp_M^*$  always holds
- (2) Symmetry:  $a \perp_M^* b \Leftrightarrow b \perp_M^* a$
- (3) Monotonicity:  $aa' \perp_M^* bb' \Rightarrow a \perp_M^* b$ .
- (4) Strong finite character: If  $a \not\perp_M^* b$ ,  
there is  $\varphi(x; b) \in \text{tp}(a/Mb)$  such that

$a' \not\perp_M^* b$  for all  $a'$  with  $F\psi(a'; b)$ .

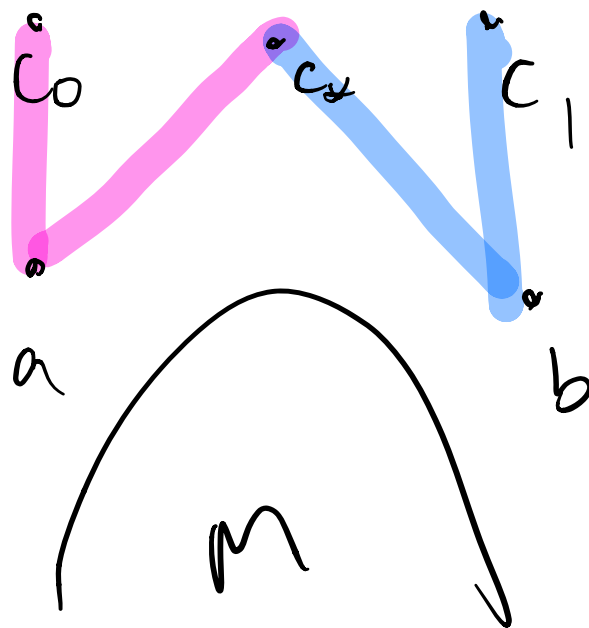
(5) The independence theorem.

If  $c_0 \equiv_M c_1$ ,  $c_0 \perp_M^* a$ ,  $c_1 \perp_M^* b$ ,  $a \perp_M^* b$ ,

then there is  $c_*$  such that

$c_* \equiv_{Ma} c_0$ ,  $c_* \equiv_{Mb} c_1$ , and  $c_* \perp_M^* ab$ .

Then,  $T$  is  $NSOP_1$ .



# The Kim-Pillay Criterion

## NSOP<sub>1</sub>

Definition Suppose  $M \neq T$ .

$a \perp_M^u b$  means  $\text{tp}(a/Mb)$

extends to a global  $M$ -finitely satisfiable type. [Equiv to  $\text{tp}(a/Ab)$  is  $M$ -fin sat].

$a \perp_M^i b$  means  $\text{tp}(a/Mb)$

extends to a global  $M$ -invariant type.

Easy exercise IF  $M \neq T$ ,

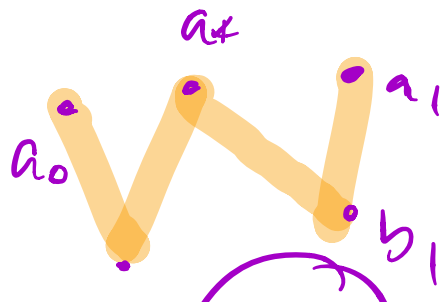
$$a \perp_M^u b \Rightarrow a \perp_M^i b \Rightarrow a \perp_M^f b \\ \Rightarrow a \perp_M^k b.$$

Definition

Define  $a \perp_M^h b \Leftrightarrow b \perp_M^u a$

and

$$a \perp_M^{ci} b \Leftrightarrow b \perp_M^i a.$$





## Proposition

$b_0$   $\left( \begin{array}{c} M \end{array} \right)$

The following are equivalent:

(1)  $T$  is  $NSOP_1$

(2) (Very weak IT for  $\perp^h$ )

If  $M \neq T$ ,  $a_0 b_0 \equiv_M a_1 b_1$  and  
 $a_i \perp^h b_i$ ,  $b_0 \perp^h b_1$ , then there

is  $a_{\neq}$  with  $a_{\neq} b_0 \equiv_M a_{\neq} b_1 \equiv_M a_0 b_0$ .

(3) (Very weak IT for  $\perp^{ci}$ )

Same as (2) but with  $\perp^{ci}$ .

=

(3)  $\Rightarrow$  (2) because  $\perp^h \Rightarrow \perp^{ci}$  hence

(1)  $\Rightarrow$  (3): Suppose (3) fails. So for some

$M \neq T$ , there are  $a_0 b_0 \equiv_M a_1 b_1$

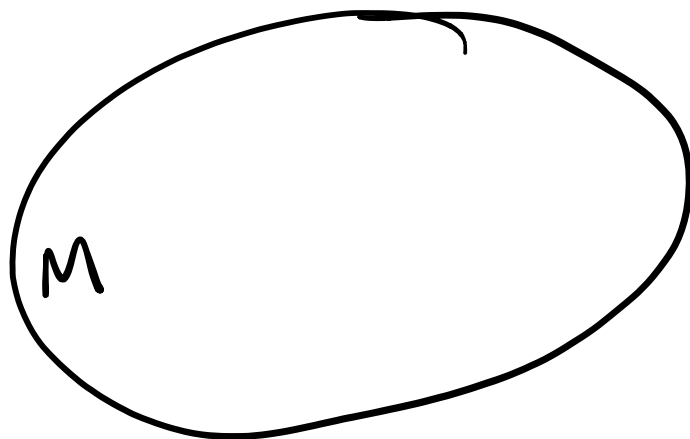
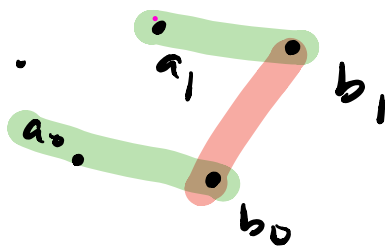
with  $a_i \perp^e b_i$  for  $i=0,1$ , and  
 $b_0 \perp^e b_1$ , but, setting

$$p(x;y) = tp(a_0, b_0 / M),$$

we have  $p(x; b_0) \cup p(x; b_1)$  is  
 inconsistent. Let  $\varphi(x;y) \in P$  be  
 a formula such that

$$\{\varphi(x; b_0), \varphi(x; b_1)\}$$

is inconsistent.



Because  $b_0 \downarrow_M^{ci} b_1$ , we know

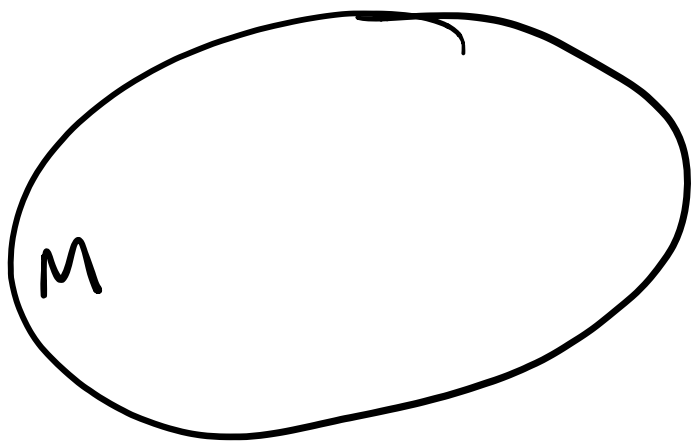
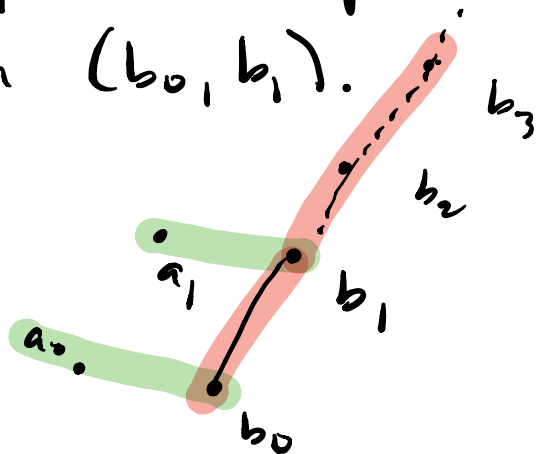
$b_1 \downarrow_M^i b_0$  so there is some global

$M$ -invariant  $q \supseteq tp(b_1/Mb_0)$ .

Note that  $b_0 \notin q/M$  so there is

a Morley sequence in  $q$  over  $M$

starting with  $(b_0, b_1)$ .



[Why? Just choose  $b_2 \neq \mathfrak{q} / M b_0 b_1$   
 $b_3 \neq \mathfrak{q} / M b_0 b_1 b_2$   
and so on.]

Then  $\{\varphi(x_i, b_i) : i < \omega\}$  is 2-inconsistent  
so  $\varphi(x_i, b_0)$  Kim-divides over  $M$ .

But we also know  $a_0 \perp_M^{c_i} b_0$ ,

so there is a global  $M$ -invariant

type  $p \supseteq tp(b_0 / M a_0)$ . Let

$(b'_i)_{i < \omega} \neq p^{\otimes \omega} / M a_0$  with  $b'_0 = b_0$ .

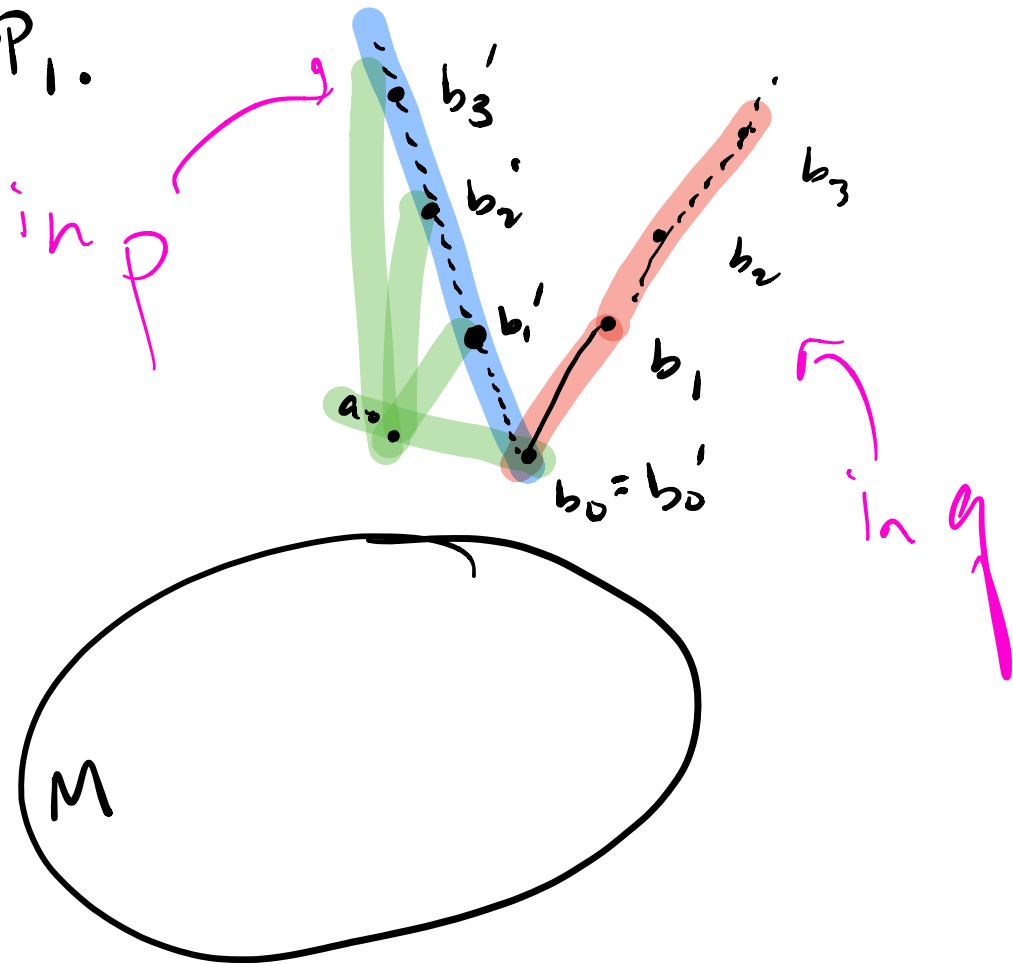
Then  $(b'_i)_{i < \omega}$  is a Morley sequence over

$M$  in  $p$ , but it is also  $M a_0$ -indiscernible

so we have  $a_0 \notin \{\psi(x_i b_i') : i \in \omega\}$ .

This contradicts Kim's lemma for  
Kim-dividing, showing  $T$  has

SOP<sub>1</sub>.



(2)  $\Rightarrow$  (1)

Suppose  $T$  has  $SOP_2$  via the formula  $\Psi(x; y)$ . Work, as before, in a Skolemized expansion  $T^{sk}$  of  $T$ . By compactness and Ramsey, there is an  $L^{sk}$ -indiscernible sequence

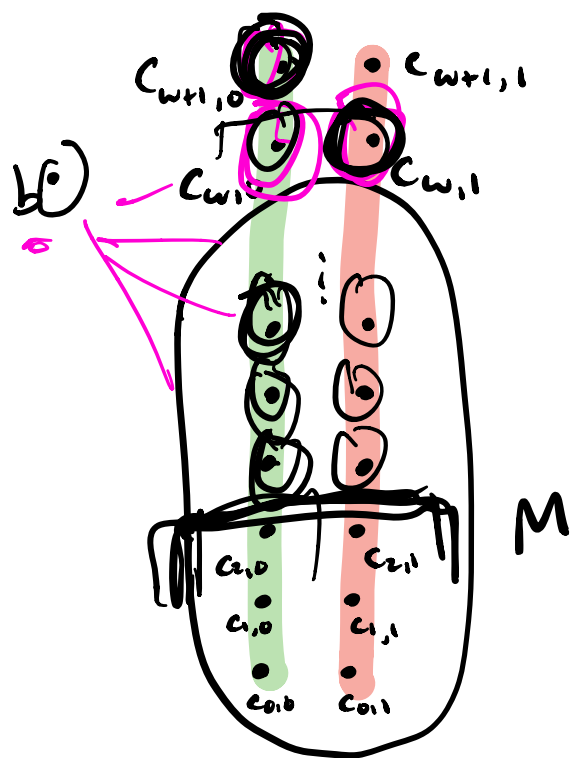
$I = (c_{i,0}, c_{i,1})_{i < \omega+2}$  such that

- $\{\Psi(x; c_{i,0}) : i < \omega+2\}$  is consistent.
- $\{\Psi(x; c_{i,1}) : i < \omega+2\}$  is 2-inconsistent.
- $c_{i,0} \equiv_{L^{sk}}^{c_{i,0} c_{i,1}} c_{i,1}$  for all  $i < \omega+2$ .

Choose  $b \models \{\Psi(x; c_{i,0}) : i < \omega+2\}$ . By Ramsey

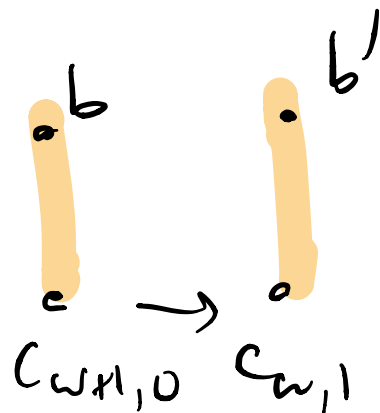
and compactness, we may assume  $I$  is  $b$ -indiscrete.

Let  $M = \text{dcl}_{\text{Sk}}(\underline{c_{w,0}}, \underline{c_{w,1}})$ .



Note that we have

$$\underline{c_{w+1,0}} \equiv_M \underline{c_{w,0}} \equiv_M \underline{c_{w,1}}$$



Choose  $b'$  such that

$$\underline{c_{w+1,0}} b \equiv_M \underline{c_{w,1}} b'$$



By indiscernibility, we have

$$\checkmark c_{w+1,0} \perp^u_M b, \text{ hence}$$

$$\checkmark c_{w+1} \perp^u_M b'; \text{ and also}$$

$$\checkmark c_{w+1} \perp^u_M c_{w+1,0}.$$

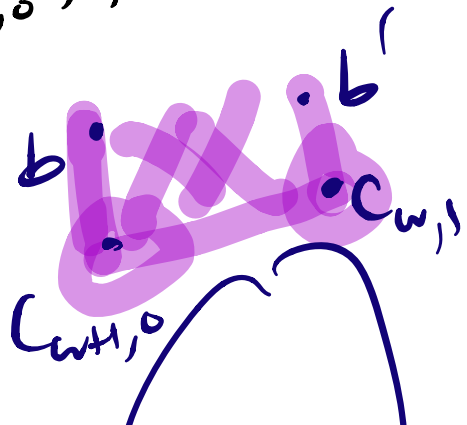
But since  $c_{w+1,0} \equiv_{c_{w+1,0} c_{w+1,1}} c_{w+1,1}$

and  $\{\varphi(x; c_{w+1}), \varphi(x'; c_{w+1,1})\}$  is

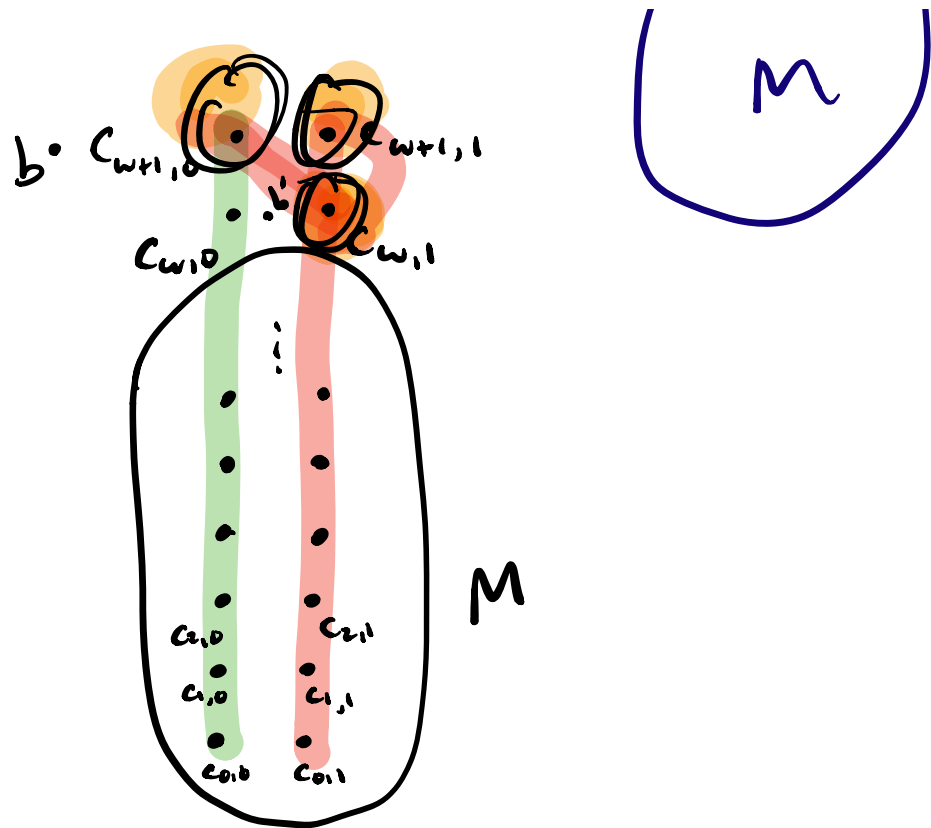
inconsistent, we have

$$\{\varphi(x; c_{w+1}), \varphi(x'; c_{w+1,0})\}$$

is inconsistent.







Since  $\varphi(x_i, c_{w+1,0}) \in \text{tp}(b/M c_{w+1,0})$

$\varphi(x_i, c_{w+1,1}) \in \text{tp}(b'/M c_{w+1,1})$ ,

we see (2) fails.  $\square$

## Proof of the Kim-Pillay theorem

First, we show that if  $\perp^*$  satisfies

(1) - (4), then for any  $M \models T$ ,

$$a \perp_M^u b \Rightarrow a \perp_M^* b.$$

Proof Suppose, towards contradiction, that

$$a \perp_M^u b + a \not\perp_M^* b. \quad \text{By strong}$$

finite character, there is some

$\varphi(x; b) \in \text{tp}(a/Mb)$  such that

$$a' \not\perp_M^* b \text{ whenever } \models \varphi(a'; b).$$

Because  $a \perp_M^u b$ , there is some

$m \in M$  such that  $\not\models \varphi(m; b)$ .

Then  $m \not\perp_M^* b$  so, by symmetry,

$b \not\perp_M^* m$  and, by monotonicity,

$b \not\perp_M^* M$ , contradicting existence.  $\square$

Now suppose  $\perp$  additionally satisfies (S).

Suppose we are given  $a_0 b_0 \equiv_M a_1 b_1$   
with  $a_0 \perp_M^h b_0$  (and hence  $a_1 \perp_M^{c_i} b_1$ )

and also  $b_0 \perp_M^h b_1$ . Then by

the claim (and symmetry), we

have  $a_0 \perp_M^* b_0$ ,  $a_1 \perp_M^* b_1$ ,  $b_0 \perp_M^* b_1$ ,

so there is, by the independence theorem,

some  $a_x$  with  $a_x b_0 \equiv_M a_x b_1 \equiv_M a_0 b_0$ .  $\square$

Teaser for next time,

Symmetry in simple theories?

Assume  $a \perp_A b$ . Then

by extn, can generate  
seq  $(a_i)_{i \in \mathbb{N}}$  with  $a_0 = a$   
 $a_i \perp_{f(A)}^{f(a)}$   $a_i \perp_A^{f(a_i)} b$

By ER, get (indisc) Morby  
seq over  $A$ ,  $Ab$ -indisc.  
Starting w/  $a$ .

Kim's lemma  $\Rightarrow b \perp^f a.$   
A

$\Leftarrow$   
 $a \downarrow^k b \not\Rightarrow a \downarrow^i b.$   
M

$\Leftarrow$   
(a)icw Modly

$a_{\geq i} \downarrow a_{\leq i}$

(a)icw  $\downarrow^k$  - Modly  
r

$a_2 a_3$   ~~$K$~~   $a_0 a_1$