

On Strong Π_1^1 -Martin-Löf Randomness

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Spector-Gandy Theorem

Theorem (Spector; Gandy)

A real $x \subseteq \omega$ is Π_1^1 if and only if it is Σ_1 over $L_{\omega_1^{\text{CK}}}$

Based on Spector-Gandy Theorem, we may define a Turing reduction in the higher setting.

Definition (Greenberg, Laurent and Monin)

A real $x \leq_{hT} y$ if there is a $\Sigma_1(L_{\omega_1^{\text{CK}}})$ relation $\Phi \subseteq 2^{<\omega} \times 2^{<\omega}$ so that

- 1** *If τ_0 and τ_1 are comparable, then for any σ_0, σ_1 , if $\Phi(\sigma_0, \tau_0)$ and $\Phi(\sigma_1, \tau_1)$, then σ_0 and σ_1 are comparable; and*
- 2** *For any $\sigma \prec x$, there is some $\tau \prec y$ such that $\Phi(\sigma, \tau)$.*

Martin-Löf Randomness

Definition (Martin-Löf)

- (i) *A Martin-Löf test is a computable collection $\{V_n : n \in \mathbb{N}\}$ of c.e. sets such that $\mu(V_n) \leq 2^{-n}$.*
- (ii) *A real y is said to pass the Martin-Löf test if $y \notin \bigcap_{n \in \omega} V_n$.*
- (iii) *A real y is Martin-Löf random if it passes all the Martin-Löf tests.*

Π_1^1 -Martin-Löf Randomness

Definition (Hjorth and Nies)

- (i) *A Martin-Löf test is a Π_1^1 collection $\{V_n : n \in \mathbb{N}\}$ of Π_1^1 -coded open sets such that $\mu(V_n) \leq 2^{-n}$.*
- (ii) *A real y is said to pass the Π_1^1 Martin-Löf test if $y \notin \bigcap_{n \in \omega} V_n$.*
- (iii) *A real y is Π_1^1 -Martin-Löf random if it passes all the Π_1^1 Martin-Löf tests.*

The collection of Π_1^1 -Martin-Löf random reals is Σ_2^0 but not Π_2^0 .

A universal Π_1^1 -Martin-Löf Random Real

There is a Π_1^1 -Martin-Löf random, left- Π_1^1 - real $\Omega_{\Pi_1^1}$ so that for any left Π_1^1 real x , $x \leq_{hT} \Omega_{\Pi_1^1}$.

Strong ML-Randomness

Definition (Kurtz)

- (i) *A generalized Martin-Löf test is a computable collection $\{V_n : n \in \mathbb{N}\}$ of c.e. sets such that $\lim_{n \rightarrow \infty} \mu(V_n) = 0$.*
- (ii) *A real y is said to pass the generalized Martin-Löf test if $y \notin \bigcap_{n \in \omega} V_n$.*
- (iii) *A real y is strong ML-random if it passes all the generalized Martin-Löf tests.*

Π_1^1 -strong ML-Randomness

Definition (Nies)

- (i) *A generalized Π_1^1 Martin-Löf test is a Π_1^1 collection $\{V_n : n \in \mathbb{N}\}$ of Π_1^1 coded open sets such that $\lim_{n \rightarrow \infty} \mu(V_n) = 0$.*
- (ii) *A real y is said to pass the generalized Π_1^1 Martin-Löf test if $y \notin \bigcap_{n \in \omega} V_n$.*
- (iii) *A real y is Π_1^1 strong ML-random if it passes all the Π_1^1 -generalized Martin-Löf tests.*

Separating Randomness Notions

Theorem (Yu; Greenberg, Laurent and Monin)

Π_1^1 strong ML-randomness is properly stronger than Π_1^1 -ML-randomness. Actually $\Omega_{\Pi_1^1}$ is not strong ML-random.

Proof.

We use an argument due to GLM. □

A technique lemma

Lemma

Suppose that $\{U_n\}_{n \in \omega}$ is a uniformly Π_1^1 -sequence of open sets. If there is a $\Sigma_1(L_{\omega_1^{\text{CK}}})$ enumeration $\{\hat{U}_{n,\gamma}\}_{n \in \omega, \gamma < \omega_1^{\text{CK}}}$ of the sequence with two numbers k and $m \geq 1$ such that for every n , $U_n = \bigcup_{\gamma < \omega_1^{\text{CK}}} \hat{U}_{n,\gamma}$ and for every $\gamma < \omega_1^{\text{CK}}$:

- (a) $\hat{U}_{n+1,\gamma} \subseteq \hat{U}_{n,\gamma}$ and each string in $\hat{U}_n$ has length at least $2^{k \cdot n}$,
- (b) $\forall \sigma \in 2^{k \cdot n - m} (\mu(\hat{U}_{n,\gamma} \cap [\sigma]) < 2^{-1+m-k \cdot n})$, and
- (c) For $\gamma < \omega_1^{\text{CK}}$ and any real z , if $z \in \hat{U}_{n,<\gamma} \setminus \hat{U}_{n,\gamma}$, where $\hat{U}_{n,<\gamma} = \bigcup_{\beta < \gamma} \hat{U}_{n,\beta}$, then $z \notin \hat{U}_{n,\beta}$ for any $\beta \geq \gamma$.

Then $\{U_n\}_{n \in \omega}$ is a generalized Π_1^1 -ML-test.

Characterizing the difference via hyperarithmetic reduction

Theorem (Chong and Yu)

For any hyperdegree $\mathbf{a}$, $\mathbf{a}$ contains a Π_1^1 -ML-random but not strong Π_1^1 -ML-random real if and only if $\mathbf{a} \geq_h \mathbf{0}'$.

Proof.

Combining the Lemma with Kučera coding. □

Π_1^1 -Difference Randomness

Definition

- (i) A d - Π_1^1 test is a Π_1^1 collection $\{V_n : n \in \mathbb{N}\}$ of Π_1^1 open sets and a Σ_1^1 -closed-set T such that $\forall n \mu(V_n \cap T) < 2^{-n}$.
- (ii) A real y is said to pass the d - Π_1^1 test if $y \notin \bigcap_{n \in \omega} V_n \cap T$.
- (iii) A real y is Π_1^1 difference random if it passes all the d - Π_1^1 tests.

Characterizing Π_1^1 -Difference-Randomness via hT -reduction

Theorem (Nies)

A Π_1^1 -ML-random real x is not Π_1^1 difference random if and only if $\Omega_{\Pi_1^1} \leq_{hT} x$.

Proof.

$\leftarrow$. A routine argument from classical randomness due to Levin-Miller-Yu.

$\rightarrow$. Fix a Σ_1^1 -closed set T and a Π_1^1 collection $\{V_n : n \in \mathbb{N}\}$ of Π_1^1 sets such that $x \in \bigcap_{n \in \omega} (T \cap V_n)$. For any k , if $m \geq k$ belongs to $\mathcal{O}$ at some stage β and $\mu(T[\gamma] \cap V_m[\beta]) \leq 2^{-k}$ for some $\gamma \geq \beta$, we put an open set with measure $\leq 2^{-k+1}$ covering the difference test at stage γ into U_k . Then $\{\bigcup_{k \geq n+1} U_k\}_{n \in \omega}$ is a Π_1^1 -ML-test and $x \notin \bigcap_n \bigcup_{k \geq n+1} U_k$. Then $\mathcal{O} \leq_{hT} x$. $\square$

Strong Π_1^1 -ML-Randomness implies Difference Randomness

The classical proof is quite easy. But the higher setting is more subtle.

Strong Π_1^1 -ML-Randomness implies Difference Randomness

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Theorem (Greenberg, Laurent and Monin)

No strong Π_1^1 -ML-Randomness is hT above $\Omega_{\Pi_1^1}$. So strong Π_1^1 -ML-randomness implies difference randomness.

Proof.

By a routine argument from classical randomness due to Levin-Miller-Yu, it can be shown that if x is strong Π_1^1 -ML-Random and $y \leq_h x$, then so is y . Then apply $\Omega_{\Pi_1^1}$. $\square$

The Borel Rank of Strong Π_1^1 -ML-Randomness

Theorem (Yu)

The collection of strong Π_1^1 -ML-random reals is not Σ_3^0 .

Proof.

By a forcing argument. The key point is to apply Π_1^1 -difference tests. □

Thanks!