

15. Martin-Lof randomness is too weak

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The notion "Martin-Lof randomness" is too weak in the sense that the halting probability Ω is Martin-Lof random, but it is left-c.e. and Turing equivalent to ϕ' , therefore it computes every Δ_2^0 set.

Kůčera - Gács theorem says Martin-Lof random sets can be arbitrary powerful as oracles.

Theorem (Kůčera - Gács)

For every set A there is a ML-random set Z s.t. $A \leq_T Z$

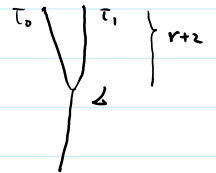
Lemma Fix measurable $S \subseteq 2^\omega$, $\delta \in 2^\omega$ and $r \in \mathbb{N}$

If $\mu(S \cap \bar{I} \delta I) \geq 2^{-(r+1)} \cdot 2^{-|\delta|}$

Then there are distinct $\tau_0, \tau_1 > \delta$ s.t. $|\tau_0| = |\tau_1| = |\delta| + r + 2$

and $\mu(S \cap \bar{I} \tau_i I) > 2^{-(r+2)} \cdot 2^{-|\tau_i|}$ for $i = 0, 1$

Proof Let $C = \{ \tau \mid |\tau| = |\delta| + r + 2 \text{ and } \tau > \delta \}$
 Note $|C| = 2^{r+2}$



Then $\mu(S \cap \bar{I} \delta I) = \sum_{\tau \in C} \mu(S \cap \bar{I} \tau I)$ (by measurability)

and it must be that $\mu(S \cap \bar{I} \tau I) \leq \mu \bar{I} \tau I = 2^{-|\tau|} = 2^{-|\delta| - (r+2)}$

As some $\mu(S \cap \bar{I} \tau_0 I) = 2^{-|\delta| - (r+2)}$ for some $\tau_0 \in C$ and $\mu(S \cap \bar{I} \tau_1 I) \leq 2^{-(r+2)} \cdot 2^{-|\tau_1|} = 2^{-|\delta| - 2(r+2)}$

$$\begin{aligned} \text{Then } \mu(S \cap \bar{I} \delta I) &\leq 2^{-|\delta| - (r+2)} + (2^{r+2} - 1) \cdot 2^{-|\delta| - 2(r+2)} \\ &= 2^{-|\delta| - (r+1)} - 2^{-|\delta| - 2(r+2)} \\ &< 2^{-|\delta| - (r+1)} \quad \square \end{aligned}$$

Proof of the theorem

Let $\{U_n\}_{n \in \mathbb{N}}$ be a universal ML-test, then U_n is a c.e. class, $\mu(U_n) \leq 2^{-n}$

and so $\mu(2^\omega - U_n) \geq 1/2$, moreover sets in $2^\omega - U_n$ are all ML-random sets.

and so $\mu(2^w - U_1) \geq 1/2$, moreover sets in $2^w - U_1$ are all $M-L$ random sets.

Let $S = 2^w - U_1$, we use sets in S to code arbitrary sets $A \in 2^w$

Let $\delta_{\tau} = \langle \tau \rangle$. then $\mu(S \cap \delta_{\tau}) = \mu(S) \cdot 2^{-|\tau|}$

Given δ_{τ} defined. Let $r = |\tau|$, IH: $\mu(S \cap \delta_{\tau}) \geq 2^{-(|\tau|+1)}$
 $= 2^{-(r+1)} \cdot 2^{-|\delta_{\tau}|}$

By the previous lemma, there are distinct $\delta_0, \delta_1 > \delta_{\tau}$ s.t.

$$|\delta_0| = |\delta_1| = |\delta_{\tau}| + r + 2 \quad \text{and}$$

$$\mu(S \cap \delta_0) > 2^{-(r+2)} \cdot 2^{-|\delta_0|}$$

Let $\delta_{\tau_0}, \delta_{\tau_1}$ be the leftmost and rightmost such strings, then

$$\begin{aligned} \mu(S \cap \delta_{\tau_0}) &> 2^{-(|\tau_0|+2)} \cdot 2^{-|\delta_{\tau_0}|} \\ &= 2^{-(|\tau|+1+1)} \cdot 2^{-|\delta_{\tau_0}|} \end{aligned}$$

Now we have defined δ_{τ} for any $\tau \in 2^{<w}$

Given $A \in 2^w$, let $Z = \bigcup_{\tau \in A} \delta_{\tau}$, clearly $Z \in S$, so Z is $M-L$ random.

Claim $A \subseteq_{\tau} Z$

To extract A from Z , Assume $A|_n = \tau$ is computed

To decide $A(n)$:

Note that for every τ , $|\delta_{\tau}| = |\tau| \cdot (|\tau| + 3) / 2$

$$[|\delta_{\langle \rangle}| = f(0) = 0, |\delta_{\tau_1}| = f(r+1) = |\delta_{\tau}| + |\tau| + 2 = f(r) + r + 2]$$

So we know $Z|_f(n) = \delta_{\tau}$, we want to decide whether

$$Z|_{f(n+1)} = \delta_{\tau_0} \text{ or } = \delta_{\tau_1} \quad \forall \epsilon.$$

if $Z|_{f(n+1)}$ is the leftmost or the rightmost δ s.t.

$$\delta > \delta_{\tau} = Z|_f(n) \quad \text{and} \quad \mu(S \cap \delta) \geq 2^{-(n+2)} \cdot 2^{-f(n+1)}$$

Key: $U_1 = 2^w - S$ is c.e. (or Σ_1^1)

So we can enumerate $\{p \mid \exists \tau \in U_1, \text{ and if for any } \theta > \delta_{\tau} \text{ } |\theta| = f(n+1)$

So we can enumerate $\{p \mid \mathbb{I}p\mathbb{I} \subset U, \gamma\}$, and if for any $\theta \succ \Delta_T$ $|\theta| = t(m)$

$$\text{if } \mu(S \cap \mathbb{I}\theta\mathbb{I}) < 2^{-(m_2)} \cdot 2^{-t(m)} \text{ i.e. } \mu(U \cap \mathbb{I}\theta\mathbb{I}) \geq 2^{-(m_2)} \cdot 2^{-t(m)}$$

we will find it in finitely many steps.

If Δ is the leftmost, in finitely many steps, we will find all θ to the left of Δ s.t. $\mu(S \cap \mathbb{I}\theta\mathbb{I}) < 2^{-(m_2)} \cdot 2^{-t(m)}$

If Δ is the rightmost, in -----

right -----

Therefore, we can decide if Δ is Δ_T or Δ_{T_1} in finitely many steps

in either case we know $A(n) = 0$ or 1 □